

Loop (with Counter)

- We can loop for ~~specific~~ specific number of times using **DJNZ** instruction

DJNZ Counter, TARGET

- This instruction will ^{→ register OR Port} Decrement the Counter then check if it ~~is~~ not equal zero then Jump to target otherwise, ~~we~~ Continue normal flow of execution

Mov R2, ~~10~~ 10 } Loop for 10 times

Loop: DJNZ R2, Loop

- 8051 Registers are 8-bit So we can loop to 255 if we need more we should use nested loop

Example: Loop for 700 times

Mov R1, ~~7~~ 7

Loop1: Mov R2, ~~100~~ 100

Loop2: DJNZ R2, Loop2 → inner

DJNZ R1, Loop1 → outer

Conditional	Jump	Jump
JZ	target	→ if $A = 0$
JNZ	target	→ if $A \neq 0$
DJNZ	R0, target	
CJNE	R0, 80 80, target	→ if $R0 \neq 80$
JC	target	→ if Carry = 1
JNC	target	→ if Carry = 0
JB	P0.1, target	→ if bit(P0.1) = 1
JNB	P1.0, target	→ if P1.0 is = 0
JBC	C, target	→ if C = 1

JBC Instruction

- This instruction check if the bit = 1
- If that is true Jump and Clear that bit

CJNE instruction (slide 222)

- This instruction is used to compare two values
- The comparison is performed by subtracting the two values without change values in registers

- It only change the Carry

CJNE A, R0, Not-EQUAL

NOT-EQUAL: ~~JNC~~ ^{→ jump not Carry} A - R0

- If the Carry = 1 that mean $[A - R0] \rightarrow$ result is negative $\rightarrow A < R0$
- If the Carry = 0 $A \geq R0$

Un Conditional Jump

Short Jump

SJMP target

- It takes only 2 Bytes
- the first byte for opcode
- the second byte is a relative

address (-128 → 127)

Address	OP Code	Instruction
0007	XX 06	SJMP target
0009	XX 00	NOP

000F target: NOP

$0F - 09 = 06$

- لاحظ أن يتغير حساب الـ relative address
بالإعتناء على عنوان الأمر التالي
المضافة إلى عنوان الذي يتغير إلى انتقال إليه

~~000F - 0009~~

$$000F - 0009 = 06$$

- لابد أن يكون الفرق في حدود 127, 128 -

- Short Jump نوع Conditional Jump

Address

0002

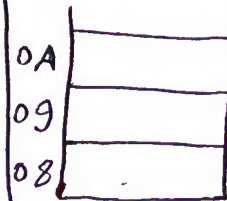
0005

Instruction

LCALL XYZ

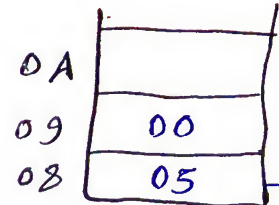
Nop

Stack



SP = 07

after LCALL



SP = 09

Low First

2 Long Jump

- It take 3 bytes
- The first byte for opcode
- The other two bytes is the absolute Address of the target

LJMP target

Look At slide 117

- pop

- push

بتنظيف علك

The CALL Instructions

1 Long Call (LCALL target)

- 3-bytes instruction
- there are two bytes for address

~~Short Call~~

2 Absolute Call (ACALL target)

- 2-bytes instruction
- 11 bits are used for address

- The CALL instruction Store the Content of Program Counter in Stack before Jump

- The RET instruction is used to POP [PC] from stack and back to instruction after call

Delay in 8051

- Delay depends on

① Crystal frequency

② Design of 8051

For 8051

* Crystal freq. = 11.0592 MHz

* Clocks per Machine Cycle

- depend on Design

- Default = 12

To Compute machine Cycle

$$\text{machine freq.} = \frac{\text{Crystal Freq.}}{\text{Clocks per Machine Cycle}}$$

$$\text{machine cycle} = \frac{1}{\text{machine freq.}}$$

Common no. of machine cycles
for Common Instruction

MOV	1
DEC	1
INC	1
DJNZ	2
LJMP	2
SJMP	2
NOP → no operation	1
MUL AB	4
PUSH	2
POP	2
RET	2

Look At Examples at
Slides 123 → 127

Example

Write Delay Subroutine to introduce
100 ms delay

$$\begin{aligned}\text{Machine freq.} &= \frac{11.0592}{12} \\ &= 921.6 \text{ KHz} \\ \text{machine cycle} &= 1/921.6 \text{ KHz} \\ &= 1.085 \mu\text{s}\end{aligned}$$

The required no. of machine cycles

$$\text{no. cycle} = \frac{100 \text{ ms}}{1.085 \mu\text{s}}$$

$$= 92165.89$$

$$\approx 92166 \text{ cycles}$$

We try to use two nested
Loop

Delay: MOV R1, ~~X~~ 200 → X

Outer: MOV R2, ~~X~~ 228 → Y

inner: DJNZ R2, inner

DJNZ R1, outer

RET

Cycle

1

1

1

1

1

MOV → 1

MOV → 1 * X

DJNZ → 1 * X * Y

DJNZ → 2 * X

RET → 2

$$92166 = 3 + 3X + 2XY$$

Assume X = 150

$$Y \approx 305 > 255$$

Assume X = 200

$$Y \approx 229$$